



**Lecture 3:
The Software Development Process**

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Organisatorisches

- ▶ Die Übung am Donnerstag, 31.10.2019, fällt aus (Reformationstag).
- ▶ Nächste Übung am Dienstag, 05.11.2019.



Where are we?

- ▶ 01: Concepts of Quality
- ▶ 02: Legal Requirements: Norms and Standards
- ▶ 03: The Software Development Process
- ▶ 04: Hazard Analysis
- ▶ 05: High-Level Design with SysML
- ▶ 06: Formal Modelling with OCL
- ▶ 07: Testing
- ▶ 08: Static Program Analysis
- ▶ 09-10: Software Verification
- ▶ 11-12: Model Checking
- ▶ 13: Conclusions



Software Development Models



Software Development Process

- ▶ A software development process is the **structure** imposed on the development of a software product.
- ▶ We classify processes according to **models** which specify
 - ▶ the artefacts of the development, such as
 - ▶ the software product itself, specifications, test documents, reports, reviews, proofs, plans etc;
 - ▶ the different stages of the development;
 - ▶ and the artefacts associated to each stage.
- ▶ Different models have a different focus:
 - ▶ Correctness, development time, flexibility.
- ▶ What does quality mean in this context?
 - ▶ What is the **output**? Just the software product, or more? (specifications, test runs, documents, proofs...)



Artefacts in the Development Process

Planning:

- Document plan
- V&V plan
- QM plan
- Test plan
- Project manual

Specifications:

- Requirements
- System specification
- Module specification
- User documents

Implementation:

- Source code
- Models
- Documentation

Possible formats:

- Documents:
 - Word documents
 - Excel sheets
 - Wiki text
 - Database (Doors)
- Models:
 - UML/SysML diagrams
 - Formal languages: Z, HOL, etc.
 - Matlab/Simulink or similar diagrams
- Source code

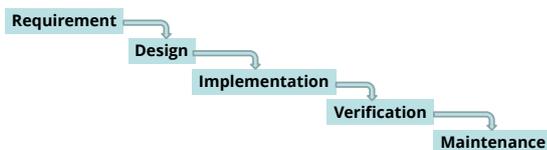
Verification & validation:

- Code review protocols
- Test cases, procedures, and test results
- Proofs



Waterfall Model (Royce 1970)

- ▶ Classical top-down sequential workflow with strictly separated phases.



- ▶ Unpractical as an actual workflow (no feedback between phases), but even the original paper did **not** really suggest this.



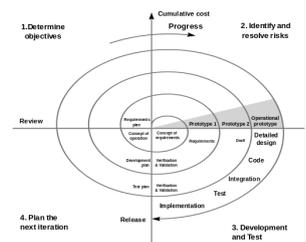
Spiral Model (Böhm 1986)

- ▶ Incremental development guided by **risk factors**

- ▶ Four phases:
 - ▶ Determine objectives
 - ▶ Analyse risks
 - ▶ Development and test
 - ▶ Review, plan next iteration

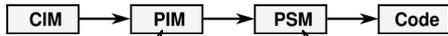
- ▶ See e.g.
 - ▶ Rational Unified Process (RUP)

- ▶ Drawbacks:
 - ▶ Risk identification is the key, and can be quite difficult



Model-Driven Development (MDD, MDE)

- Describe problems on abstract level using a *modeling language* (often a *domain-specific language*), and derive implementation by model transformation or run-time interpretation.
- Often used with UML (or its DSLs, eg. SysML)



- Variety of tools:
 - Rational tool chain: Enterprise Architect, Rhapsody, Artisan
 - Studio, Platform-independent model (link/StateWork)
 - EMF (Eclipse Modeling Framework), Platform-specific model
- Strictly sequential development
- Drawbacks: high initial investment, limited, reverse engineering and change management (code changes to model changes) is complex

* Proprietary DSL – not related to UML



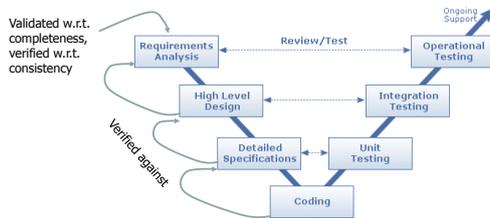
Agile Methods

- Prototype-driven development
 - E.g. Rapid Application Development
 - Development as a sequence of prototypes
 - Ever-changing safety and security requirements
- Agile programming
 - E.g. Scrum, extreme programming
 - Development guided by functional requirements
 - Process structured by rules of conduct for developers
 - Rules capture best practice
 - Less support for non-functional requirements
- Test-driven development
 - Tests as *executable specifications*: write tests first
 - Often used together with the other two

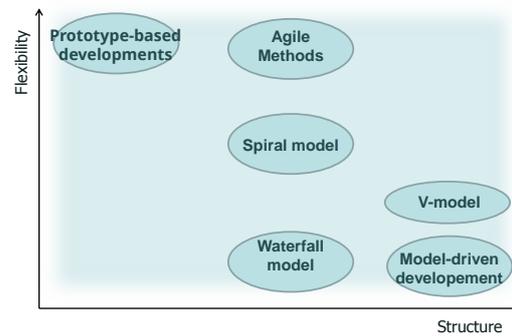


V-Model

- Evolution of the waterfall model:
 - Each phase supported by corresponding verification & validation phase
 - Feedback between next and previous phase
- Standard model for public projects in Germany
 - ... but also a general term for models of this „shape“
- Current: V-Modell XT („extreme tailoring“)
 - Shape gives dependencies, not development sequence



Software Development Models



from S. Paulus: Sichere Software



Development Models for Safety-Critical Systems

Development Models for Critical Systems

- Ensuring safety/security needs structure.
 - ...but *too much* structure makes developments bureaucratic, which is *in itself* a safety risk.
 - Cautionary tale: Ariane-5
- Standards put emphasis on **process**.
 - Everything needs to be planned and documented.
 - Key issues: **auditability, accountability, traceability**.
- Best suited development models are variations of the V-model or spiral model.
- A new trend? V-Model XT allows variations of original V-model, e.g.:
 - V-Model for initial developments of a new product
 - Agile models (e.g. Scrum) for maintenance and product extensions



Auditability and Accountability

- Version control and configuration management is **mandatory** in safety-critical development (auditability).
- Keeping track of all artifacts contributing to a particular instance (**build**) of the system (**configuration**), and their **versions**.
- Repository** keeps all artifacts in all versions.
 - Centralised: one repository vs. distributed (every developer keeps own repository)
 - General model: check out – modify – commit
 - Concurrency: enforced **lock**, or **merge** after commit.
- Well-known systems:
 - Commercial: ClearCase, Perforce, Bitkeeper...
 - Open Source: Subversion (centralised); Git, Mercurial (distributed)



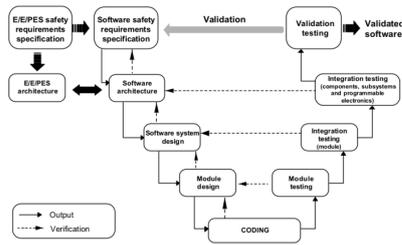
Traceability

- The idea of being able to follow requirements (in particular, safety requirements) from requirement spec to the code (and possibly back).
- On the simplest level, an Excel sheet with (manual) links to the program.
- More sophisticated tools include DOORS:
 - Decompose requirements, hierarchical requirements
 - Two-way traceability: from code, test cases, test procedures, and test results back to requirements
 - E.g. DO-178B requires all code derives from requirements
- The SysML modelling language has traceability support:
 - Each model element can be traced to a requirement.
 - Special associations to express traceability relations.



Development Model in IEC 61508

- IEC 61508 in principle allows any development model, but:
 - It requires safety-directed activities in each phase of the life cycle (safety life cycle, cf. last lecture).
 - Development is one part of the life cycle.
- The only development model mentioned is a V-model:

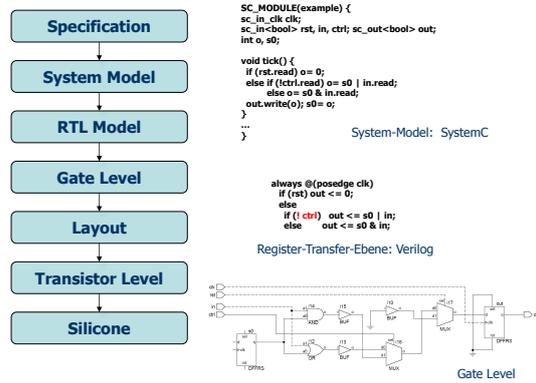


Development Model in DO-178B/C

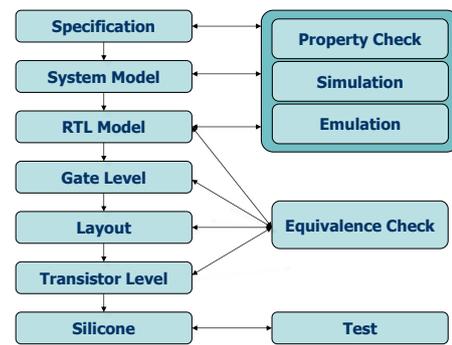
- DO-178B/C defines different *processes* in the SW life cycle:
 - Planning process
 - Development process, structured in turn into
 - Requirements process
 - Design process
 - Coding process
 - Integration process
 - Verification process
 - Quality assurance process
 - Configuration management process
 - Certification liaison process
- There is no conspicuous diagram, but the Development Process has sub-processes suggesting the phases found in the V-model as well.
 - Implicit recommendation of the V-model.



Development Model for Hardware



Development Model for Hardware



Basic Notions of Formal Software Development

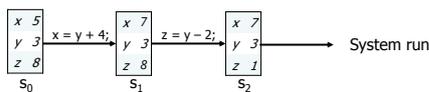
Formal Software Development

- In a formal development, properties are stated in a rigorous way with a **precise mathematical semantics**.
- Formal specification requirements can be **proven**.
- Advantages:**
 - Errors can be found early in the development process.
 - High degree of confidence into the system.
 - Recommend use of formal methods for high SILs/EALs.
- Drawbacks:**
 - Requires a lot of effort and is thus expensive.
 - Requires qualified personnel (that would be *you*).
- There are tools which can help us by
 - finding (simple) proofs for us (model checkers), or
 - checking our (more complicated) proofs (theorem provers).



Formal Semantics

- States** and transitions between them:



- Operational semantics** describes relation between states and transitions:

$$\frac{s \vdash e \rightarrow n}{s \vdash x = e \rightarrow s[x/n]} \quad \text{hence:} \quad \frac{s_0 \vdash y + 4 \rightarrow 7}{s_0 \vdash x = y + 4 \rightarrow s_1}$$

- Formal proofs**, e.g. proving

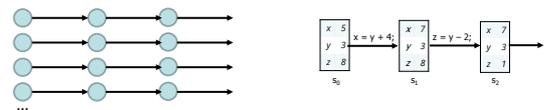
$$x = y + 4; z = y - 2;$$

yields the same final state as

$$z = y - 2; x = y + 4;$$

Semantics of Programs and Requirements

- Set of all possible system runs



- Requirements** related to safety and security:

- Requirements on single states ?
- Requirements on system runs ?
- Requirements on sets of system runs ?

Alpern & Schneider
Clarkson & Schne



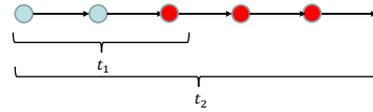
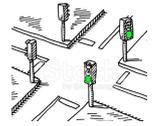
Some Notions

- ▶ Let b, t be two traces then
 $b \leq t$ iff. $\exists t'. t = b \cdot t'$ i.e. b is a *finite prefix* of t
- ▶ A **property** is a set of infinite execution traces (like a program)
 - ▶ Trace t satisfies property P , written $t \models P$, iff $t \in P$
- ▶ A **hyperproperty** is a set of sets of infinite execution traces (like a set of programs)
 - ▶ A system (set of traces) S satisfies H iff $S \in H$
 - ▶ An observation Obs is a finite set of finite traces
 - ▶ $Obs \leq S$ (Obs is a prefix of S) iff
 Obs is an observation and $\forall m \in Obs. \exists t \in S. m \leq t$



Requirements on States: Safety Properties

- ▶ Safety property S : „Nothing bad happens“
 - ▶ i.e. the system will never enter a *bad* state
 - ▶ E.g. „Lights of crossing streets do not go green at the same time“
- ▶ A bad state:
 - ▶ can be **immediately** recognized;
 - ▶ **cannot be sanitized** by following states.
- ▶ S is a safety property iff
 $\forall t. t \notin S \Rightarrow (\exists t_1. t_1 \leq t \Rightarrow \forall t_2. t_1 \leq t_2 \Rightarrow t_2 \notin S), t_1$ finite



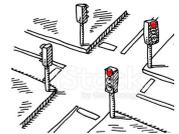
Proving Safety Properties

- ▶ In the previous specification, t_1 is **finite**. As a consequence,
 - ▶ a property is a safety property if and only if its violation can be detected on a finite trace.
- ▶ Safety properties are typically proven by induction
 - ▶ Base case: initial states are good (= not bad)
 - ▶ Step case: each transition transforms a good state again in a good state
- ▶ Safety properties can be enforced by run-time monitors
 - ▶ Monitor checks following state in advance and allows execution only if it is a good state



Requirements on Runs: Liveness Properties

- ▶ Liveness property L :
 - ▶ „Good things will happen eventually“
 - ▶ E.g. „my traffic light will go green eventually“
- ▶ A good thing is always possible and possibly infinite.
- ▶ L is a liveness property iff
 - ▶ $\forall t. \text{finite}(t) \rightarrow \exists t_1. t \cdot t_1 \in L$
 - ▶ i.e. all finite traces t can be extended to a trace in L .



* Achtung: „eventually“ bedeutet „irgendwann“ oder „schlussendlich“ aber **nicht** „eventuell“ !



Satisfying Liveness Properties

- ▶ Liveness properties cannot (!) be enforced by run-time monitors.
- ▶ Liveness properties are typically proven by the help of well-founded orderings
 - ▶ Measure function m on states s
 - ▶ Each transition decreases m
 - ▶ $t \in L$ if we reach a state with minimal m
- ▶ E.g. measure denotes the number of transitions for the light to go green



Requirements on Sets of Runs: Safety Hyperproperties

- ▶ Safety hyperproperty: „System never behaves bad“
 - ▶ No bad thing happens in a finite set of finite traces
 - ▶ (the prefixes of) different system runs do not exclude each other
 - ▶ E.g. „the traffic light cycle is always the same“
- ▶ A bad system can be recognized by a bad observation (set of finite runs)
 - ▶ A bad observation cannot be sanitized regards less how we continue it or add additional system runs
 - ▶ E.g. two system runs having different traffic light cycles
- ▶ S is a safety hyperproperty iff (see [safety property](#)):



$$\forall T. T \notin S \Rightarrow (\exists Obs. Obs \leq T \Rightarrow \forall T'. Obs \leq T' \Rightarrow T' \notin S)$$



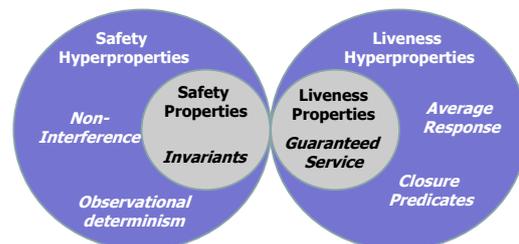
Requirements on Sets of Runs: Liveness Hyperproperties

- ▶ Liveness hyperproperty S :
 „The system will eventually develop to a good system“
 - ▶ Considering any finite part of a system behavior, the system eventually develops into a „good“ system (by continuing appropriately the system runs or adding new system runs)
 - ▶ E.g. „Green light for pedestrians can always be omitted“
- ▶ L is liveness hyperproperty iff
 $\forall T. \exists G. T \leq G \wedge G \in L$
 - ▶ T is a finite set of finite traces (observation)
 - ▶ Each observation can be explained by a system G satisfying L
- ▶ Examples:
 - ▶ Average response time
 - ▶ Closure operations in information flow control
 - ▶ Fair scheduling



Landscape of (Hyper)Properties

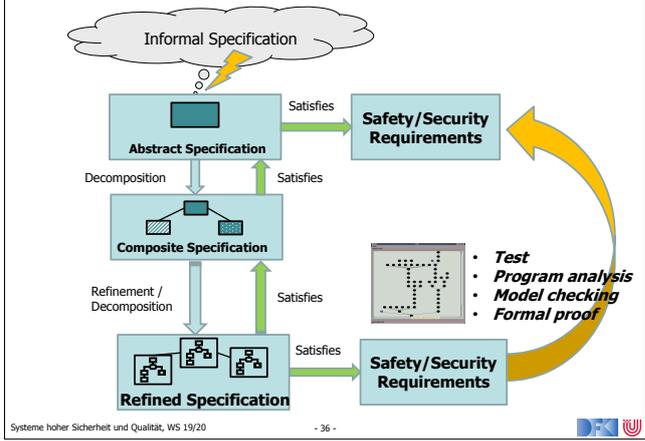
- ▶ Each (hyper-) property can be represented as a combination of safety and liveness (hyper-) properties.



Structuring the Formal Development



The Global Picture



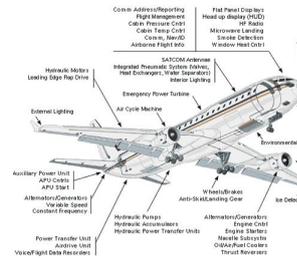
Structuring the Development

- ▶ Horizontal structuring:
 - ▶ Modularization into components
 - ▶ Composition and Decomposition
 - ▶ Aggregation
- ▶ Vertical structuring:
 - ▶ Abstraction and refinement from design specification to implementation
 - ▶ Declarative vs. imperative specification
 - ▶ Inheritance of properties
- ▶ Views:
 - ▶ Addresses multiple aspects of a system
 - ▶ Behavioral model, performance model, structural model, analysis model (e.g. UML, SysML)

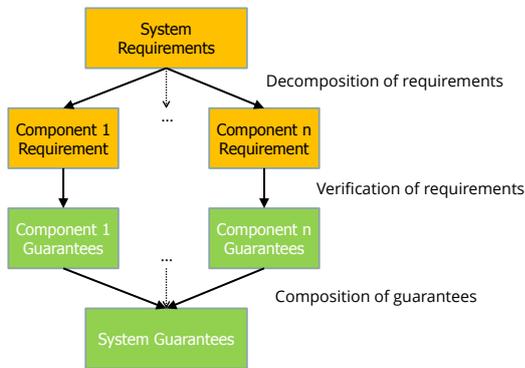


Horizontal Structuring (informal)

- ▶ Composition of components
 - ▶ Dependent on the individual layer of abstraction
 - ▶ E.g. modules, procedures, functions,...
- ▶ Example:

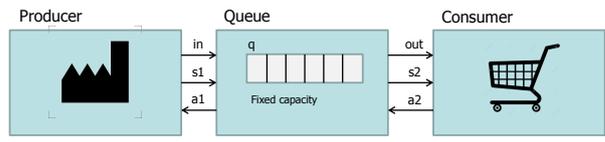


Modular Structuring of Requirements



Mutual Dependencies: Assume/Guarantee

- ▶ Safety requirement: Queue does not lose any items.



```

Loop:
if (s1 == a1) {
  send(x, in); s1 = not s1;
}

Loop:
if (s1 != a1 && |q| < max) {
  enq(q, in);
  a1 = not a1;
}
if (s2 == a2 && |q| > 0) {
  deq(q, out);
  s2 = not s2;
}

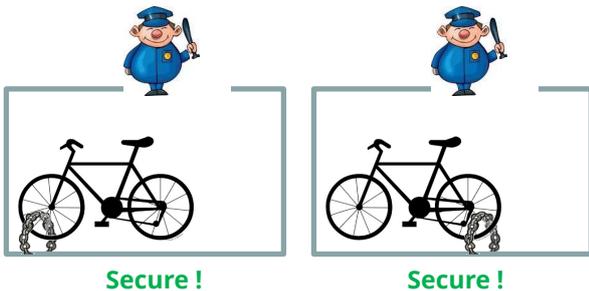
Loop:
if (s2 != a2) then {
  read(y, out);
  a2 = not a2;
  consume(y);
}
    
```

- ▶ Components depend on each other!
- ▶ Initialization ?



Composition of Security Guarantees

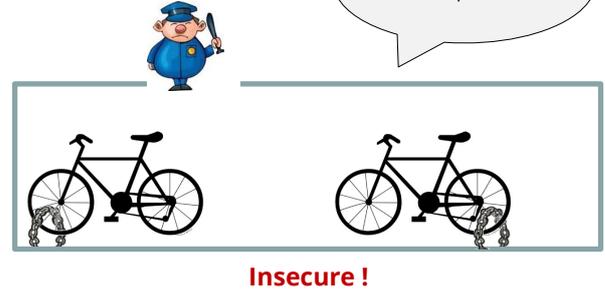
Only complete bicycles are allowed to pass the gate.



Composition of Security Guarantees

Only complete bicycles are a

Security properties are non-compositional !



Concurrent shared variable programs are non-compositional

```

long long x;

Thread1() {
  x = 1;
}
// @post: x == 1

Thread2() {
  x = (1 << 64);
}
// @post: x == (1 << 64)

(Thread1() || Thread2());
// @post: x == 1 or x == (1 << 64)

```

Global variable

Post conditions hold in absence of concurrent threads

Does composition hold?

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Concurrent shared variable programs are non-compositional

```

long long x;

(Thread1() || Thread2());

// @post: x == 1 or x == (1 << 64) or x == (1 << 64) + 1

```

- This post-condition cannot be derived from any logical composition of the original post-conditions of Thread1() and Thread2()
- For writing a 128bit integer to memory, two writes on the memory bus are required. As a consequence, the final value of x may also be (1 << 64) + 1

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Vertical Structuring - Refinement

Idea: start at an abstract description and add step by step details

From abstract specification to an implementation

- What do we want to refine?
 - Algorithm: algebraic refinement
 - Data: data refinement
 - Process: process refinement
 - Events: action refinement

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Algebraic Refinement

Stack

```

empty: stack;
pop(stack):stack;
push(int, stack):stack

pop(empty) = empty
pop(push(x, y)) = y

```

Refinement

Implementing stacks by lists

```

empty ↦ []
push ↦ (:)
pop ↦ safetail

```

Satisfies

List

```

[] :: [a]
head :: [a]-> a
(:) :: a-> [a]-> [a]
tailSafe :: [a]-> [a]

tailSafe xs = if null xs then [] else tail xs

```

To prove: safetail([]) == []
safetail(y:xs) == y

Refinement preserves properties of stack by transitivity of the logic!

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Even More Refinements

- Data refinement
 - Abstract datatype is „implemented“ in terms of the more concrete datatype
 - Simple example: define stack with lists
- Process refinement
 - Process is refined by excluding certain runs
 - Refinement as a reduction of underspecification by eliminating possible behaviours
- Action refinement
 - Action is refined by a sequence of actions
 - E.g. a stub for a procedure is refined to an executable procedure

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Conclusion & Summary

- Software development models: structure vs. flexibility
- Safety standards such as IEC 61508, DO-178B suggest development according to V-model.
 - Specification and implementation linked by verification and validation.
 - Variety of artefacts produced at each stage, which have to be subjected to external review.
- Safety / Security Requirements
 - Properties: sets of traces
 - Hyperproperties: sets of properties
- Structuring of the development:
 - Horizontal – e.g. composition
 - Vertical – refinement (e.g. algebraic, data, process...)

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