

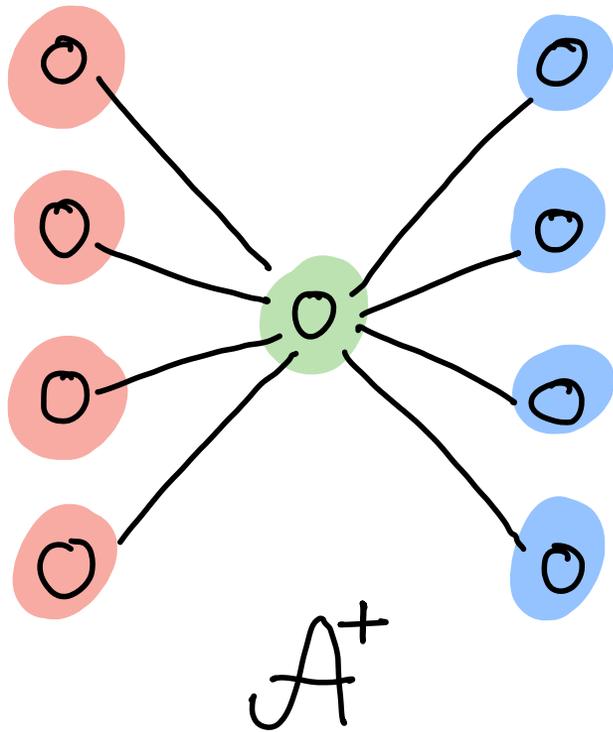


# Indiscernible Sequences in

monadically stable and monadically NIP  
classes

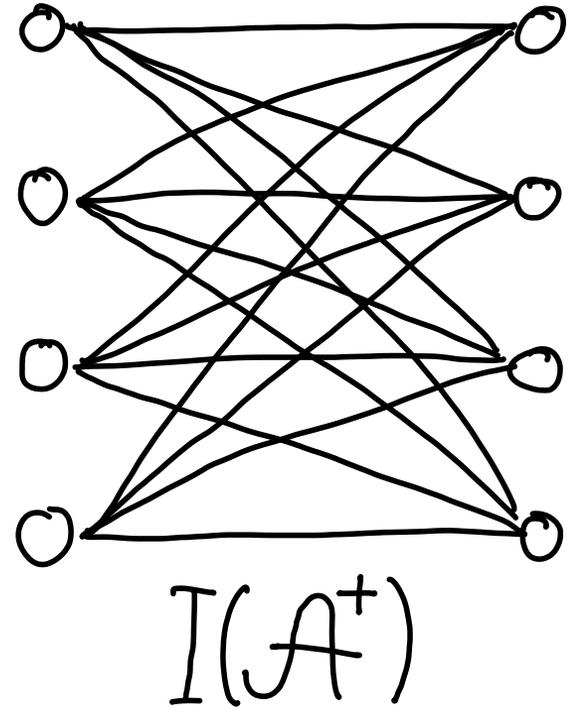
Jan Dreier, *Nikolas Mählmann*, Amer Mouawad,  
Sebastian Siebertz, Alexandre Vigny

# (simple) Interpretations



$$\overline{I} = (\mathcal{V}, \mathcal{S})$$

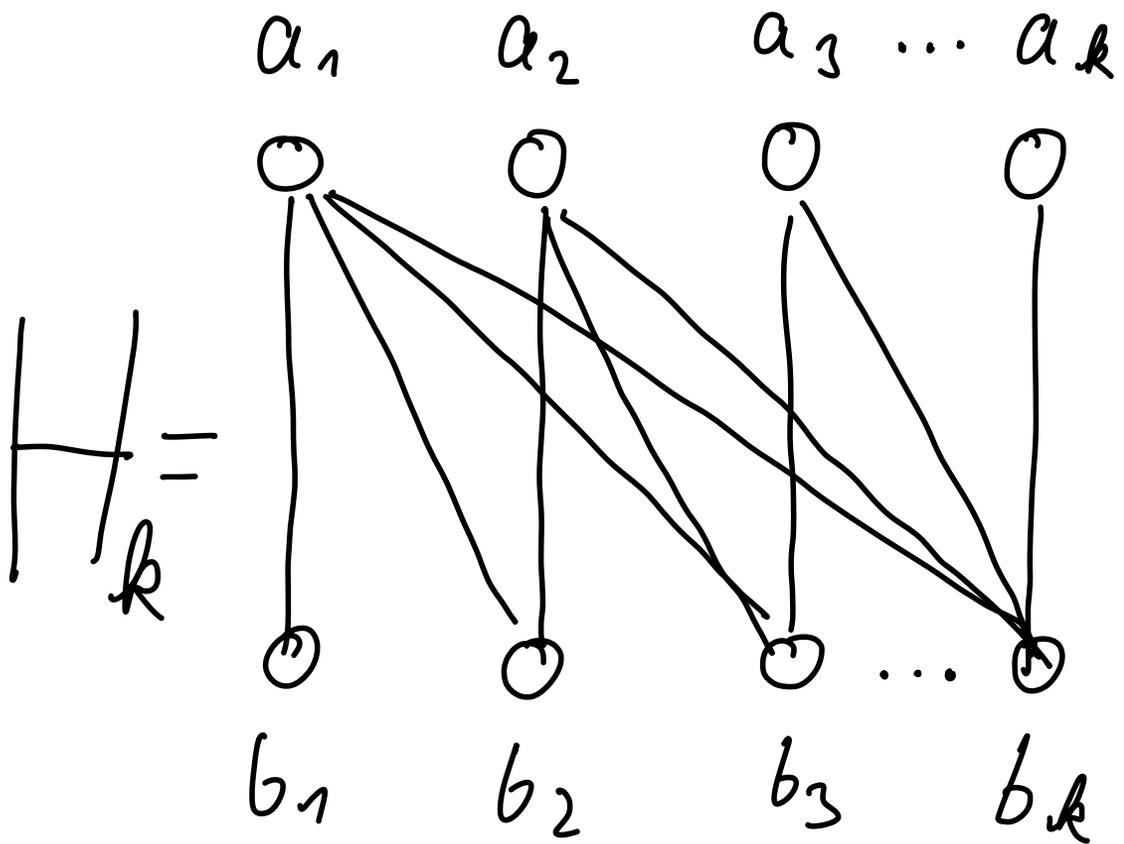
→



$$f(x, y) = \text{red } x \wedge \text{blue } y \wedge \text{dist} \leq 2(x, y)$$
$$\mathcal{S}(x) = \neg \text{green } x$$

## (monadic) Stability:

A class  $\mathcal{C}$  of structures is (monadically) stable if no interpretation interprets the class of all half-graphs in (a coloring) of  $\mathcal{C}$ .



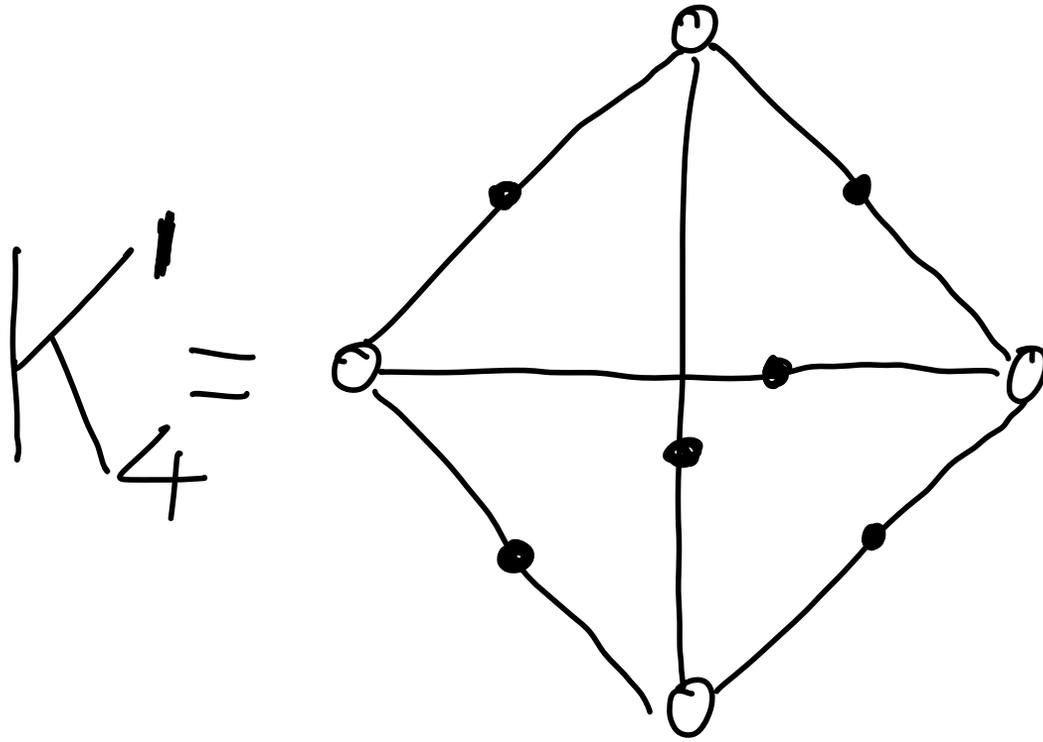
$$a_i \sim b_j \iff i \leq j$$

$\mathcal{C}$  is mon. stable iff.

$$\forall I \exists k \forall A^+ \in \mathcal{C}^+ : H_k \notin I(A^+)$$

## (monadic) NIP

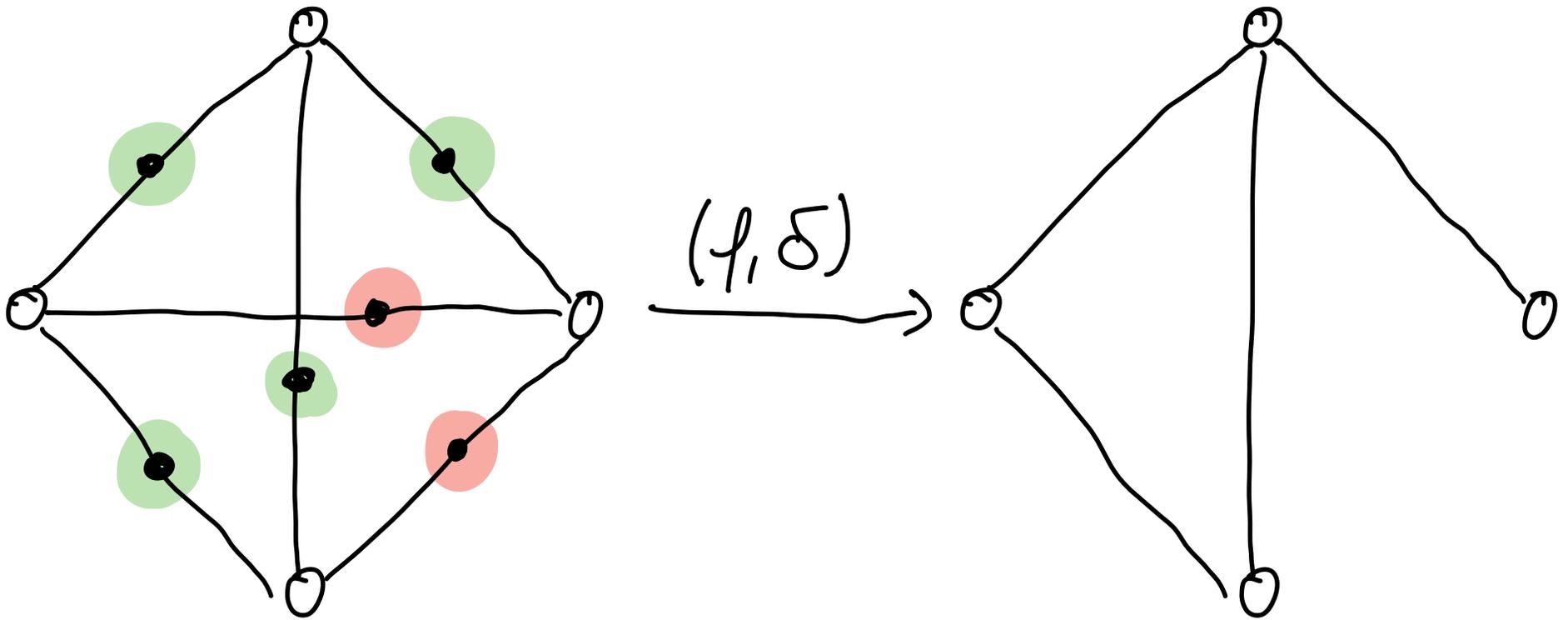
A class  $\mathcal{C}$  of structures is (monadically) NIP if no interpretation interprets the class of all graphs in (a coloring) of  $\mathcal{C}$ .



$\mathcal{C}$  is mon. NIP iff.

$$\forall I \exists k \forall A^+ \in \mathcal{C}^+ : K'_k \notin I(A^+)$$

# Encoding arbitrary graphs in $K'_k$



$$f(x, y) = \exists z: z \wedge x \sim z \sim y$$

$$\delta(x) = \neg z \wedge \neg x$$

# Overview of selected graph classes

mon. stable

U~~X~~

$\subseteq$

mon. NIP

U~~X~~

nowhere dense

U~~X~~

$\not\subseteq$   
 $\not\supseteq$

bd. twinwidth

U~~X~~

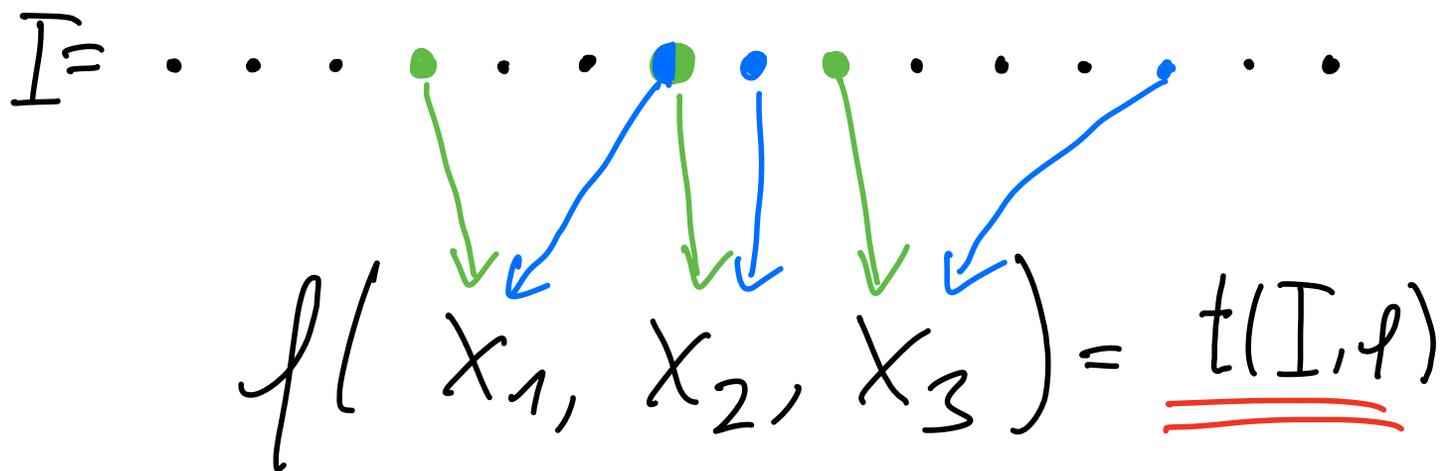
bd. treewidth

$\subseteq$

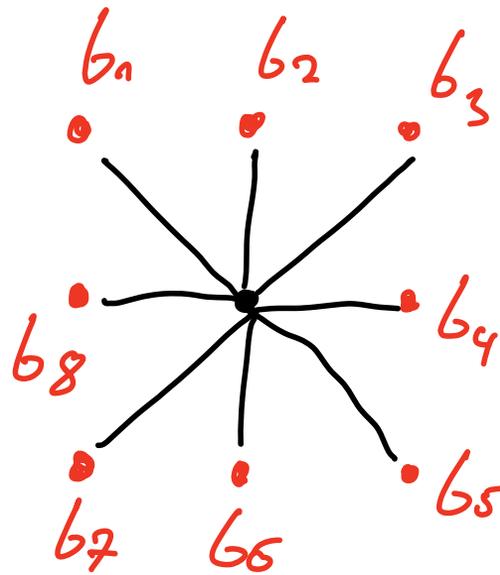
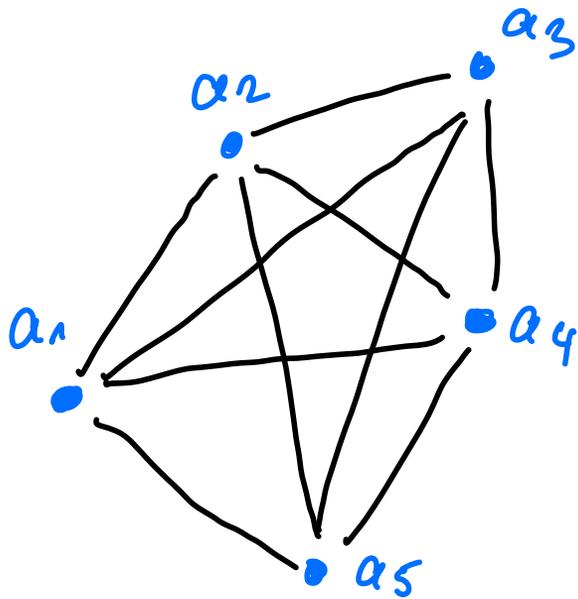
bd. cliquewidth

# Indiscernible Sequences

$I = (a_1, \dots, a_n)$  is a  $\Delta$ -indiscernible sequence, if every  $f(x_1, \dots, x_k) \in \Delta$  has a constant truth value on every  $k$ -element subsequence of  $I$ .



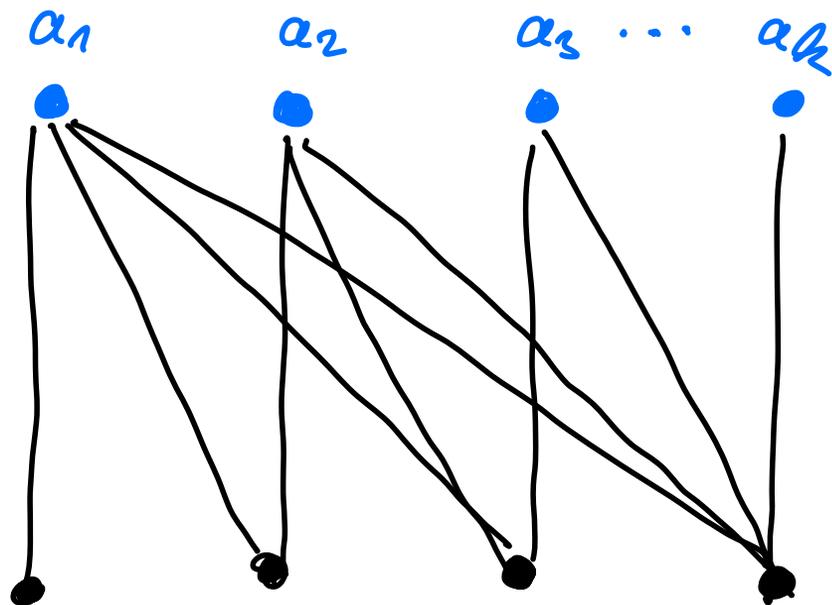
# Example I



Both  $(a_1, \dots, a_5)$  and  $(b_1, \dots, b_8)$  are  $\{\beta\}$ -indiscernible sequences for:

$$f(x_1, x_2) = x_1 \sim x_2$$

# Example II



$(a_1, \dots, a_k)$  is a  $\{\beta\}$ -indiscernible sequence for:

$$f(x_1, x_2) = \exists y: x_1 \sim y \wedge x_2 \not\sim y$$

# Properties of Indiscernible Sequences

In general for every fixed set  $\Delta$

- Every large enough sequence contains a  $\Delta$ -ISeq.

In stable classes

- $\Delta$ -ISeqs have polynomial size

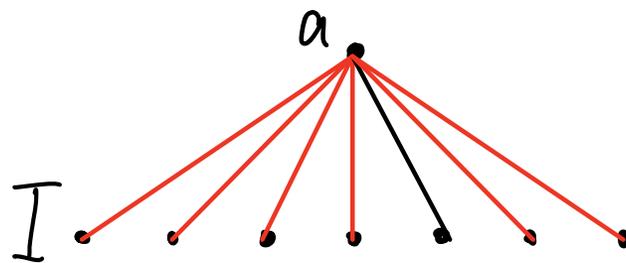
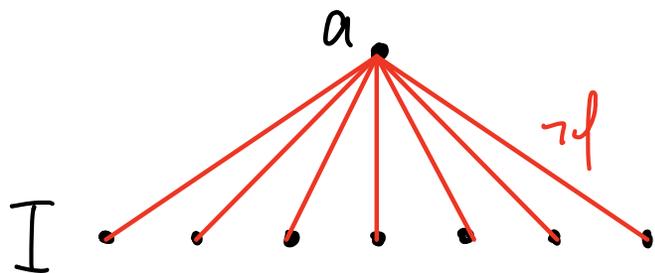
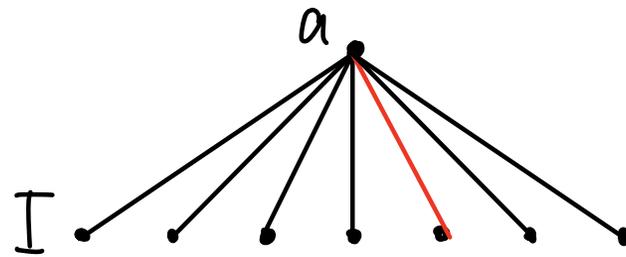
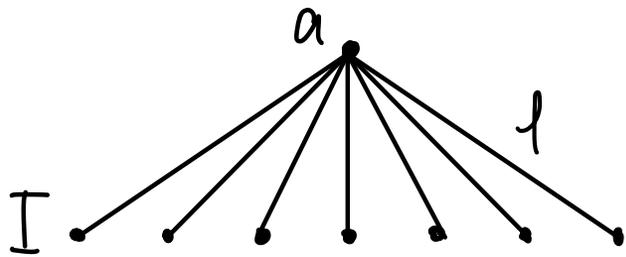
- $\Delta$ -ISeqs are **totally indiscernible**

i.e. remain indiscernible when permuted

Up next: Our Results!

A class  $\mathcal{C}$  is mon. stable  
iff.

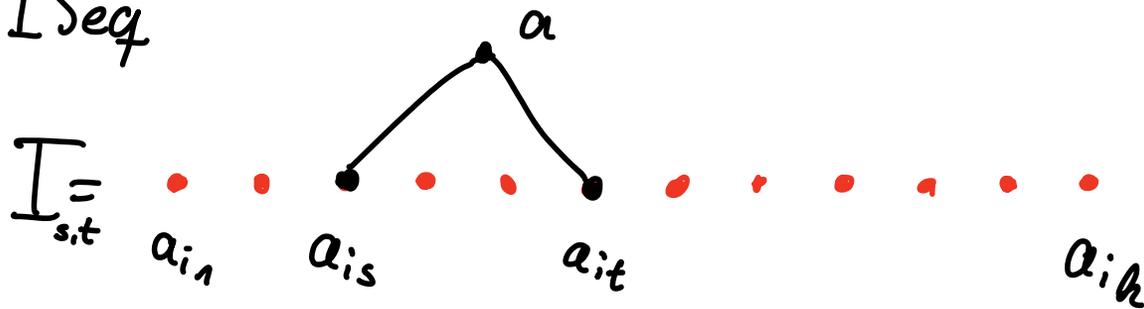
for every coloring  $\mathcal{C}^+$  and formula  $\varphi(x,y)$  there exists  
a set  $\Delta$  s.t. for every  $\Delta$ -ISeg  $I$  in a structure  $\mathcal{A}^+ \in \mathcal{C}^+$   
for every element  $a \in \mathcal{A}^+$  one of the following cases applies:



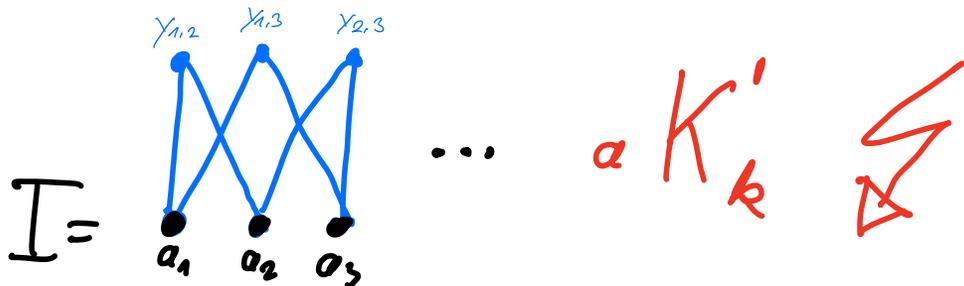
# Proof Sketch



by permuting elements of  $I$  we find for every  $s \neq t \in [k] \times [k]$  an  $I$ Seq



witnessing  $f_{s,t}(x_1, \dots, x_k) = \exists y : [y \sim x_i] \leftrightarrow [i \in \{s, t\}]$



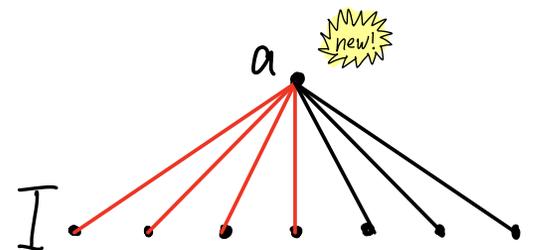
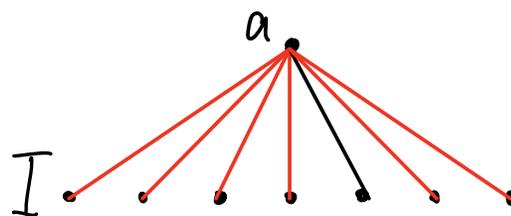
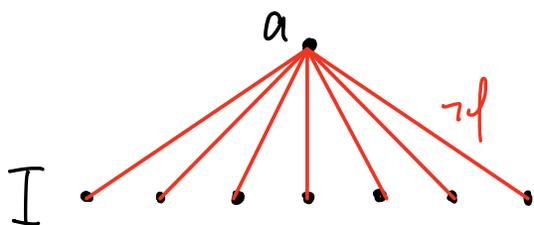
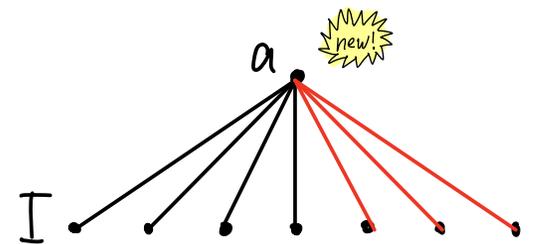
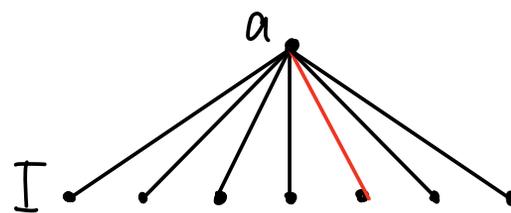
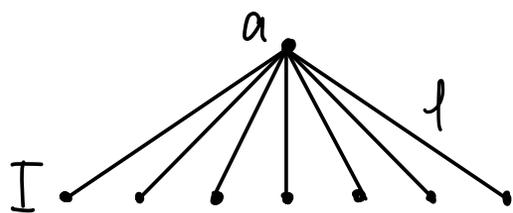
A class  $\mathcal{C}$  is mon. NIP

iff.

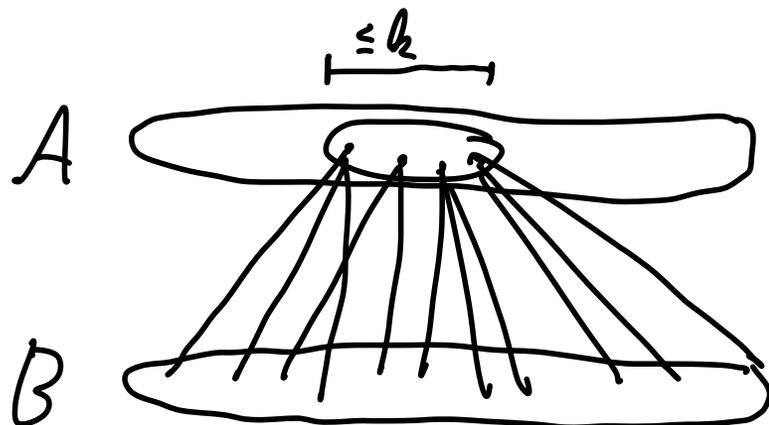
for every coloring  $\mathcal{C}^+$  and formula  $\varphi(x,y)$  there exists

a set  $\Delta$  s.t. for every  $\Delta$ -ISeq  $I$  in a structure  $\mathcal{A}^+ \in \mathcal{C}^+$

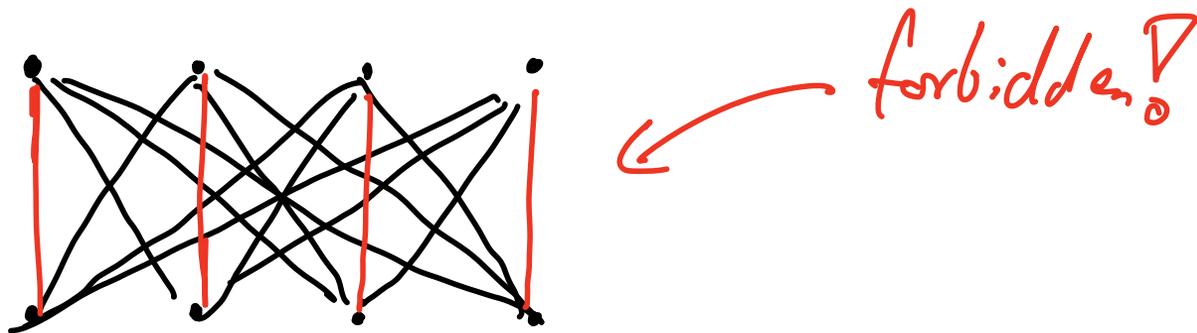
for every element  $a \in \mathcal{A}^+$  one of the following cases applies:



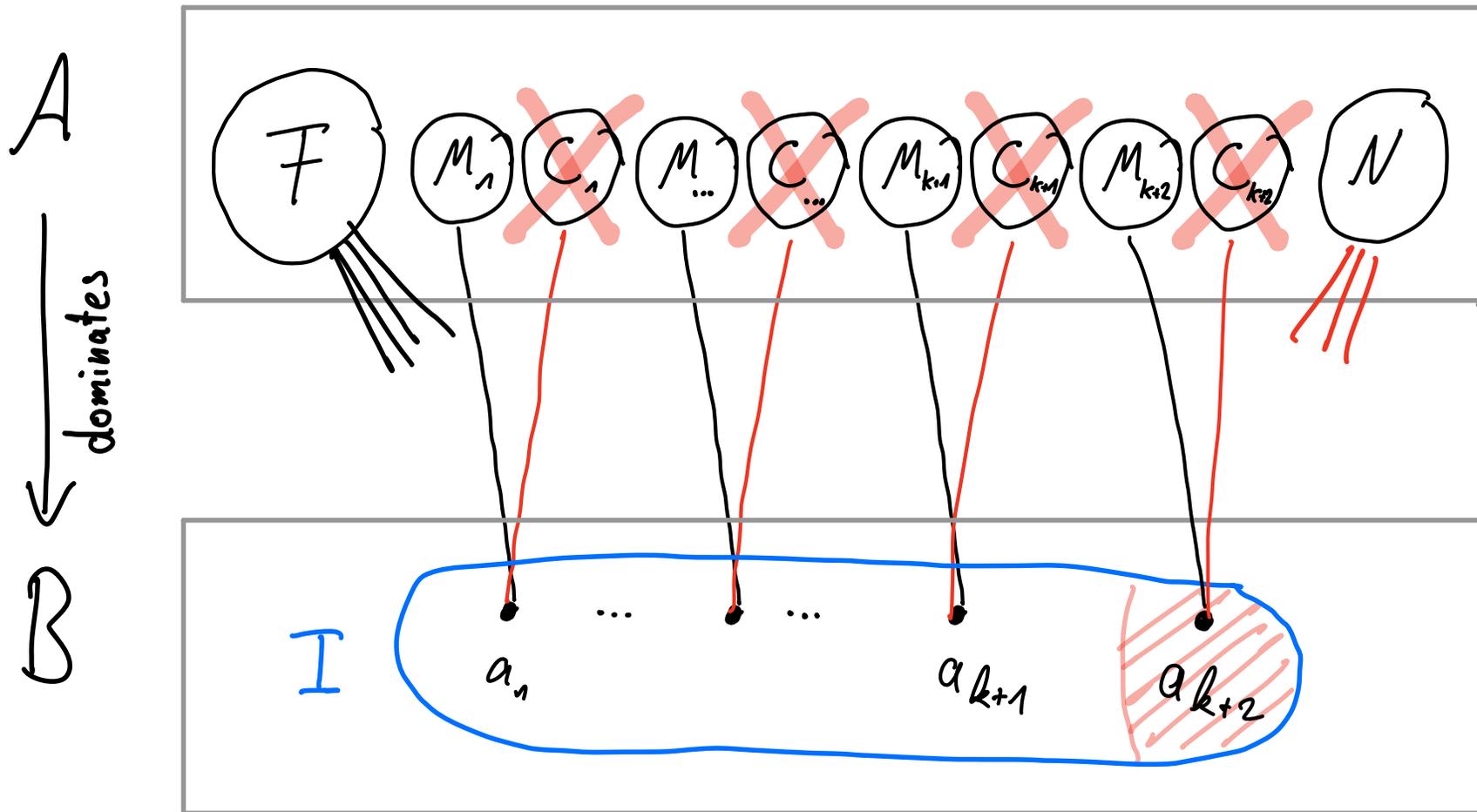
# $k$ - Dominating Sets in powers of nowhere dense graphs



powers of nowhere dense graphs are mon. stable  
and exclude a co-matching



# $k$ - Dominating Sets in powers of nowhere dense graphs



→ a polynomial kernel

# Further Applications

- improved bounds for Ramsey Numbers

→ in mon. stab.  $K_s$ -free classes we find independent sets of size

$$\Omega(|G|^{\frac{1}{s-1} + \delta})$$

- a regularity lemma for mon. stable classes

- polynomial kernels for powers of nowhere dense graphs

→ Independent Set

→ Dominating Set

- ???

Thank you for listening! ▽

Are there questions?

# Further Results

