

DFK

Systeme hoher Sicherheit und Qualität
Universität Bremen, WS 2017/2018

Lecture 06:



Formal Modeling with OCL

Christoph Lüth, Dieter Hutter, Jan Peleska

mit Folien v. Bernhard Beckert (KIT)

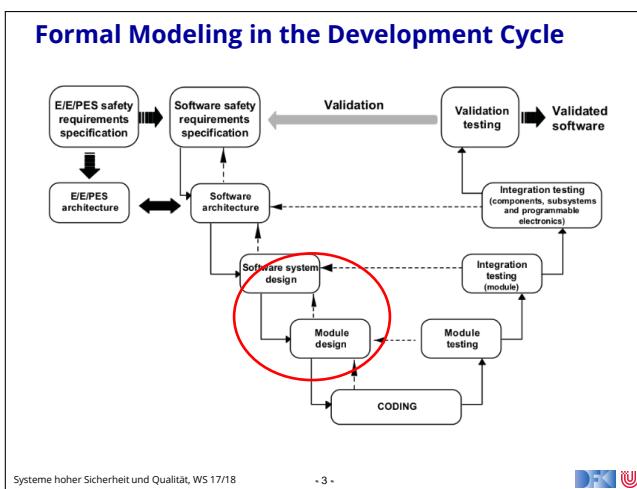
Universität Bremen

Where are we?

- ▶ 01: Concepts of Quality
- ▶ 02: Legal Requirements: Norms and Standards
- ▶ 03: The Software Development Process
- ▶ 04: Hazard Analysis
- ▶ 05: High-Level Design with SysML
- ▶ 06: Formal Modelling with OCL
- ▶ 07: Testing
- ▶ 08: Static Program Analysis
- ▶ 09-10: Software Verification
- ▶ 11-12: Model Checking
- ▶ 13: Conclusions

Systeme hoher Sicherheit und Qualität, WS 17/18

+ 2 +



What is OCL?

- ▶ OCL is the **Object Constraint Language**.
- ▶ What is OCL?
 - ▶ „A formal language used to describe expressions on UML models. These expressions typically specify invariant conditions that must hold for the system being modeled or queries over objects described in a model.“ (OCL standard, §7)
- ▶ Why OCL?
 - ▶ „A UML diagram, such as a class diagram, is typically not refined enough to provide all the relevant aspects of a specification. There is, among other things, a need to describe additional constraints about the objects in the model.“ (OCL standard, §7.1)

Systeme hoher Sicherheit und Qualität, WS 17/18

- 4 -



Characteristics of the OCL

- ▶ OCL is a pure **specification language**
 - ▶ OCL expressions do not have side effects
- ▶ OCL is **not** a programming language.
 - ▶ Expressions are not executable (though some may be)
- ▶ OCL is **typed** language
 - ▶ Each expression has type; all expressions must be well-typed
 - ▶ Types are classes, defined by class diagrams

Systeme hoher Sicherheit und Qualität, WS 17/18

- 5 -

DFK

Usage of the OCL

- ▶ as a query language
- ▶ to specify **invariants** on classes and types in the class
- ▶ to specify type invariant for Stereotypes
- ▶ to describe pre- and post conditions on Operations and Methods
- ▶ to describe guards
- ▶ to specify target (sets) for messages and actions
- ▶ to specify constraints on operations
- ▶ to specify derivation rules for attributes for any expression over a UML model.

(OCL standard, §7.1.1)

Systeme hoher Sicherheit und Qualität, WS 17/18

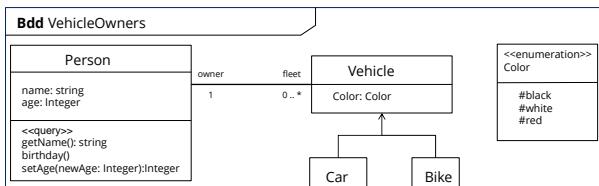
- 6 -



OCL by Example

Universität Bremen

Why is SysML not enough?



What about requirements like:

- ▶ The minimal age of car owners
- ▶ The maximal number of cars (of a specific color) owned
- ▶ The maximal number of owners of a car

Systeme hoher Sicherheit und Qualität, WS 17/18

- 8 -



OCL Basics

► The language is **typed**: each expression has a type.

► Multiple-valued logic (true, false, undefined).

► Expressions always live in a **context**:

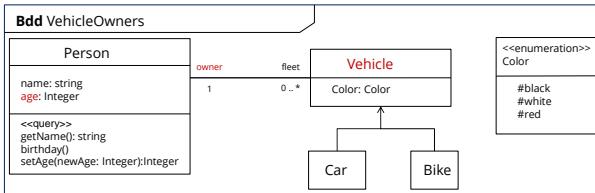
► **Invariants** on classes, interfaces, types.

```
context Class
  inv Name: expr
```

► **Pre/postconditions** on operations or methods

```
context Type :: op(al: Type, ..., an: Type) : Type
  pre Name: expr
  post Name: expr
```

Invariants of Classes



"A vehicle owner must be at least 18 years old"

```
context Vehicle
inv: self.owner.age >= 18
```

Basic types and operations

► **Integer (\mathbb{Z})**

OCL-Std. §11.5.2

► **Real (\mathbb{R})**

OCL-Std. §11.5.1

- Integer is a subclass of Real
- round, floor from Real to Integer

► **String (Zeichenketten)**

OCL-Std. §11.5.3

- substring, toReal, toInteger, characters, etc.
- Relationen auf Real, Integer, String

► **Boolean (Wahrheitswerte)**

OCL-Std. §11.5.4

- or, xor, and, implies
- Relationen auf Real, Integer, String

Collection Types

Sequence, Bag, OrderedSet, Set

OCL-Std. §11.6, §11.7

► Operations on all collections:

- size, includes, count, isEmpty, flatten
- Collections are always „flattened“

► **Set**

- union, intersection

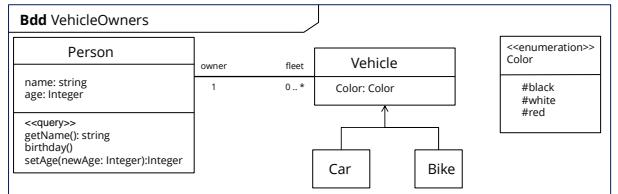
► **Bag**

- union, intersection, count

► **Sequence**

- first, last, reverse, prepend, append

Collections



"Nobody has more than 3 vehicles"

```
context Person
Inv: self.fleet->size <= 3
```

Collection Types: Quantification

We can quantify over collections:

OCL-Std. §11.9.1

► **Universal quantification**:

```
coll->forAll(elem: Type | expr[elem]) : Boolean
```

► **Existential quantification**:

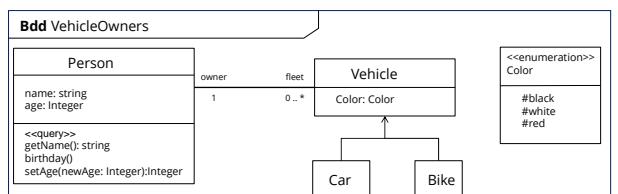
```
coll->exists(elem: Type | expr[elem]) : Boolean
```

► **Comprehension operator**:

```
coll->select(elem: Type | expr[elem]) : Coll[Type]
```

where **expr** is an expression of type Boolean.

Universal Quantification



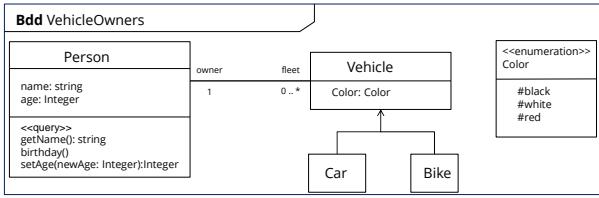
"All vehicles of a person are black"

```
context Person
inv: self.fleet->forAll(v | v.color = #black)
```

"No person has more than three black vehicles"

```
context Person
inv: self.fleet->select(v | v.color = #black)->size <= 3
```

Universal Quantification



"A person younger than 18 owns no cars"

```

context Person
inv: self.age < 18 implies
    self.fleet -> forAll(v | not v.ocllsKindOf(Car))

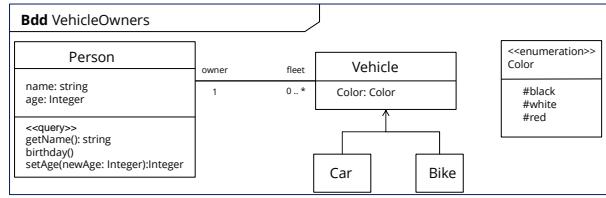
```

Systeme hoher Sicherheit und Qualität, WS 17/18

- 17 -



Existential Quantification



"There is a red car"

```

context Car
inv: Car.allInstances() -> exists(c | c.color = #red)

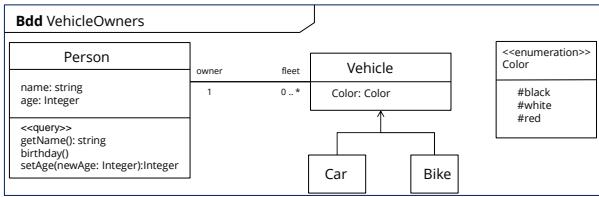
```

Systeme hoher Sicherheit und Qualität, WS 17/18

- 18 -



Pre/Post Conditions



"If **setAge(a)** is called with a non-negative argument **a**, then **a** becomes the new value of the attribute **age**."

```

context Person::setAge(a:int)
pre: a >= 0
post: self.age = a

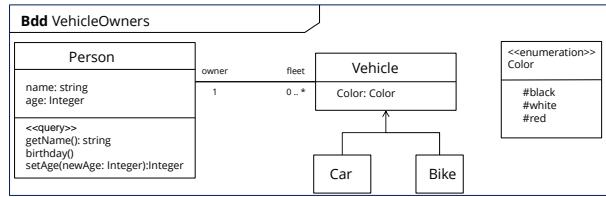
```

Systeme hoher Sicherheit und Qualität, WS 17/18

- 19 -



Pre/Post Conditions



"Calling **birthday()** increments the age of a person by 1."

```

context Person::birthday()
post: self.age = self.age@pre + 1

```

Systeme hoher Sicherheit und Qualität, WS 17/18

- 20 -



Modelling Dynamic Aspects

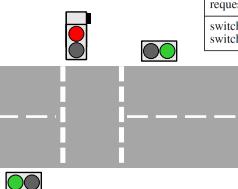
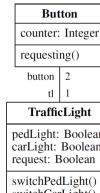
- Block diagrams model the **static structure** of the system: classes, attributes and the type of the operations. The possible **system states** are all instances of these model types.
- Invariants and pre/post conditions can be used to model the **dynamic aspects** of the system. In particular, they model all possible **state transitions** between the system states.
- An operation can become **active** (there is a state transition emanating from it) if the invariant holds, and the precondition holds. If there are no active state transitions, the system is **deadlocked**.
 - Deadlocks must be avoided.

Systeme hoher Sicherheit und Qualität, WS 17/18

- 22 -



Example: The Traffic Light



Systeme hoher Sicherheit und Qualität, WS 17/18

- 23 -



Example: The Traffic Light



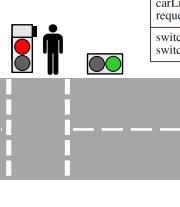
```

context requesting()
pre: tl.pedLight = false
post: tl.request = true
post: counter = counter@pre + 1

context switchPedLight()
pre: request = true
post: pedLight != pedLight@pre
post: request = false

context switchCarLight()
post: carLight != carLight@pre
inv: not(pedLight = true and
carLight = true)

```



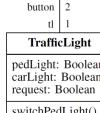
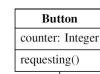
pedLight: False
carLight: True
request: False
counter: 0

Systeme hoher Sicherheit und Qualität, WS 17/18

- 24 -



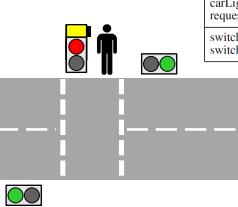
Example: The Traffic Light



```

pedLight: False
carLight: True
request: True
counter: 1

```

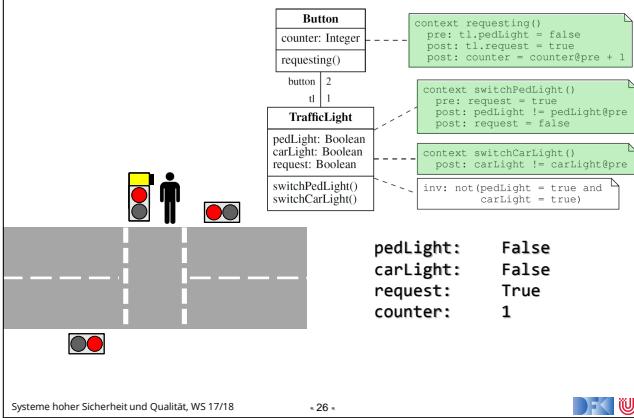


Systeme hoher Sicherheit und Qualität, WS 17/18

- 25 -



Example: The Traffic Light

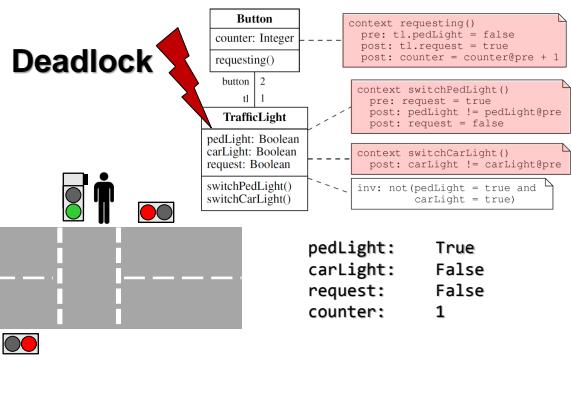


Systeme hoher Sicherheit und Qualität, WS 17/18

- 26 -



Example: The Traffic Light



Systeme hoher Sicherheit und Qualität, WS 17/18

- 27 -



OCL Details

Universität Bremen

Model types

- ▶ Model types are given by
 - ▶ Attributes,
 - ▶ Operations, and
 - ▶ Associations of the model
- ▶ Navigation along the association
 - ▶ If cardinality is 1, type is of target type \mathbb{T}
 - ▶ Otherwise, it is $\text{Set}(\mathbb{T})$
- ▶ User-defined operations in expressions have to be stateless (stereotype $\langle\langle \text{query} \rangle\rangle$)

Systeme hoher Sicherheit und Qualität, WS 17/18

- 29 -



Collection Types: Iterators

- ▶ Quantifiers are a special case of iterators.
 - ▶ Think of all/any in Haskell defined via foldr
- ▶ All iterators defined via iterate OCL-Std. §7.6.6


```
coll->iterate(elem: Type, acc: T = initial_expr
| expr[elem, acc] : Coll[T]
```

where `expr` of type \mathbb{T} denotes a function on `elem` and `acc`

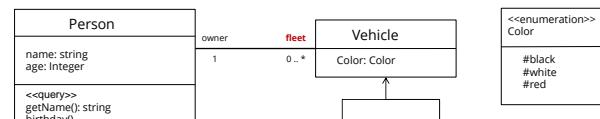
```
c.iterate(e: T, acc: T = v) =
  acc = v;
  for (Enumeration e = c.elements(); e.hasMoreElements();) {
    acc = expr[e, acc];
    e = e.nextElement();
  }
  return acc;
```

Systeme hoher Sicherheit und Qualität, WS 17/18

- 30 -



Collection Types: Iterators



"A person owns at most 3 black vehicles"

```
context Person
inv: self.fleet->iterate(v; acc: Integer = 0
| if (v.color = #black)
  then acc + 1 else acc
endif ) <= 3
```

Systeme hoher Sicherheit und Qualität, WS 17/18

- 31 -



Undefinedness in OCL

- ▶ Each domain of a basic type has two values denoting "**undefinedness**": OCL-Std §A.2.1.1
 - ▶ `null` or \perp stands for "undefined", e.g. if an attribute value has not been set or is not defined (Type `OclVoid`)
 - ▶ `invalid` or \top stands for "invalid" and signals an error in the evaluation of an expression (e.g. division by 0, or application of a partial function) (Type `OclInvalid`)
 - ▶ As subtypes: $OclInvalid \subseteq OclVoid \subseteq$ all other types
- ▶ Undefinedness is **propagated**.
 - ▶ In other words, all operations are **strict**: „an invalid or null operand causes an invalid result“.

Systeme hoher Sicherheit und Qualität, WS 17/18

- 32 -



The OCL Logic

- ▶ Exceptions to strictness:
 - ▶ Boolean operators (see below)
 - ▶ Case distinction
 - ▶ Test on definedness: `oclIsUndefined` with

$$\text{oclIsUndefined}(e) = \begin{cases} \text{true} & \text{if } e = \perp \vee e = \text{null} \\ \text{false} & \text{otherwise} \end{cases}$$
- ▶ The domain type for `Boolean` also contains `null` and `invalid`.
 - ▶ The resulting logic is **four-valued**.
 - ▶ It is a **Kleene-Logic**: $A \rightarrow B \equiv \neg A \vee B$
 - ▶ Boolean operators (`and`, `or`, `implies`, `xor`) are **non-strict on both sides**.
 - ▶ But equality (like all other relations) is strict: $\perp = \perp$ is \perp

Systeme hoher Sicherheit und Qualität, WS 17/18

- 33 -



OCL Boolean Operators: Truth Table

b_1	b_2	$b_1 \text{ and } b_2$	$b_1 \text{ or } b_2$	$b_1 \text{ xor } b_2$	$b_1 \text{ implies } b_2$	$\text{not } b_1$
false	false	false	false	false	true	true
false	true	false	true	true	true	true
true	false	false	true	true	false	false
true	true	true	true	false	true	false
false	ε	false	ε	ε	true	true
true	ε	ε	true	ε	ε	false
false	\perp	false	\perp	\perp	true	true
true	\perp	\perp	true	\perp	\perp	false
ε	false	false	ε	ε	ε	ε
ε	true	ε	true	ε	true	ε
ε	ε	ε	ε	ε	ε	ε
ε	\perp	\perp	\perp	\perp	\perp	ε
\perp	false	false	\perp	\perp	\perp	\perp
\perp	true	\perp	true	\perp	true	\perp
\perp	\perp or ε	\perp	\perp	\perp	\perp	\perp

► Legend: \perp is invalid, ε is null.

OCL-Std §A .2.1.3, Table A.2

OCL Style Guide

- Avoid **complex** navigation („Loose coupling“).
 - Otherwise changes in models break OCL constraints.
- Always choose **adequate context**.
- „Use of **allInstances()** is **discouraged**“
- Split up invariants if possible.
- Consider defining **auxiliary operations** if expressions become too complex.

Summary

- OCL is a typed, state-free specification language which allows us to denote constraints on models.
- We can define or models much more precise.
 - Ideally: no more natural language needed.
- OCL is part of the more „academic“ side of UML/SysML.
 - Tool support is not great, some tools ignore OCL, most tools at least type-check OCL, hardly any do proofs.
- However, in critical system development, the kind of specification that OCL allows is **essential**.
- Try yourself: USE – Tool <http://useocl.sourceforge.net>
 Martin Gogolla, Fabian Büttner, and Mark Richters. [USE: A UML-Based Specification Environment for Validating UML and OCL](#), Science of Computer Programming, 69:27-34, 2007.